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# The Self-Diagnosability of a Computer

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**Abstract**—Maximum capability for self-diagnosis with minimum additional hardware is the goal of every designer of a general purpose computer today. A yardstick with which the self-diagnosability of a system can be measured is proposed.

A self-diagnosable computer can be described as a system consisting of two interconnected but independent machines: the main processor  $M_1$  and a much smaller machine  $M_2$  (about 5 to 10 percent of the size of  $M_1$ ), which is capable of (programmatically) detecting and locating a fault in  $M_1$ . This fault location should be pinpointed within a small number of replaceable modules (integrated circuit chips, parallel-plate packages, or printed circuit cards) [1], [2].

The most commonly employed technique for diagnosis is to prepare a list of a complete set of tests  $T = \{T_1, T_2, \dots, T_n\}$  such that every failure in the system will cause one or more of these tests to fail [2]–[4]. Let the set  $F = \{F_1, F_2, \dots, F_m\}$  represent all possible single failure cases in the system. By taking the intersection of the sets of suspects for the failing test cases  $T_{i1}, T_{i2}, \dots, T_{ir}$  one arrives at a fault  $F_i$ . Let  $k_1, k_2, \dots, k_m$  be the number of suspected modules under the faults  $F_1, F_2, \dots, F_m$ , respectively. In other words, during the Maintenance Routine [3] run, if tests  $T_{i1}, T_{i2}, \dots, T_{ir}$  fail, and the rest of the tests pass, then from a look-up table we arrive at the conclusion that fault  $F_i$  has occurred, and in order to correct this fault  $F_i$  we have to either replace  $k_i$  number of modules or examine each of these  $k_i$  modules by some other means and replace the bad one.

Clearly then, if  $N$  = total number of modules used in the machine,

$$\sum_{i=1}^m k_i \geq N. \quad (1)$$

Let  $p_i$  be the probability of occurrence of failure  $F_i$ , for  $i=1, 2, \dots, m$ . Then assuming that at a given instance exactly one fault has occurred,

$$\sum_{i=1}^m p_i = 1. \quad (2)$$

Number  $R_i = 1/k_i$  is an indicator of the efficiency with which fault  $F_i$  can be repaired. The diagnostic efficiency of the entire system can be represented by

$$R = \frac{1}{\sum_{i=1}^m k_i p_i}. \quad (3)$$

This number  $R$  can be called the "resolution" of the entire system. The comparative figure of merit of a diagnostic subsystem is then

$$\frac{R}{\text{cost}} \quad (4)$$

where the cost includes the cost of hardware in  $M_2$ , of software, of development, and of running time of the maintenance routine.

If the maintenance routine only detects, and does not locate a fault, then  $R$  assumes its minimum possible value

$$R_{\min} = \frac{1}{N}. \quad (5)$$

This implies that one has to examine all  $N$  modules of the machine to locate the faulty module.

The resolution  $R$  is maximum when every failure can be traced down exactly to one module, i.e.  $k_i = 1$ , for  $1 \leq i \leq n$ . Then from (3), resolution becomes

$$R_{\max} = \frac{1}{\sum_{i=1}^m p_i} = 1. \quad (6)$$

If all modules are assumed to have equal probability of failure, then the probability of occurrence of failure  $F_i$  is given by

$$p_i = \frac{k_i}{\sum_{i=1}^m k_i} \quad (7)$$

and the resolution of the machine by substituting (7) in (3) turns out to be

$$R_e = \frac{\sum_{i=1}^m k_i}{\sum_{i=1}^m k_i^2}. \quad (8)$$

In absence of any statistical data available on the probability of various failures, (8) would be a good index of the diagnosability of a system.

In the author's opinion the resolution in (8) is a very important figure in the specification of any machine with diagnostic capability. The manufacturer should specify it, and the customer should ask for it. As discussed above, in general,  $R_e$  will have a value between 1 and  $1/N$ .

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